

Mathematics for Computer Science

Self-evaluation exercises for Chapter 3

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Exercise 3.1 (cf. Problem 3.4)

Describe a simple procedure which, given a positive integer argument, n , produces a width n array of truth-values whose rows would be all the possible truth-value assignments for n propositional variables. For example, for $n = 2$, the array would be:

T	T
T	F
F	T
F	F

Your description can be in English, or a simple program in some familiar language such as Python or Java. If you do write a program, be sure to include some sample output.

Exercise 3.2 (from the classroom test of 3 October 2018)

Find a disjunctive normal form for the following formula:

$$(P \text{ or } Q) \text{ implies not}(R \text{ and } P)$$

Use either a truth table, or logical equivalences.

Exercise 3.3 (cf. Problem 1.16(c)-(d))

1. Let P be a propositional formula. Explain why P is valid if and only if $\text{not}(P)$ is *not* satisfiable.

2. A finite set of propositional formulas $X = \{P_1, \dots, P_n\}$ is *consistent* if there exists an assignment of truth values to all the variables which appear in any formulas in which all propositions are true. For example, the following set is consistent:

$$\{P \text{ and not}(Q), Q \text{ or } R\},$$

but the following one is not:

$$\{A \text{ and not}(A)\}.$$

Construct a formula S such that S is valid if and only if X is *not* consistent.

Exercise 3.4 (cf. Problem 3.18(c))

We have seen during the classroom exercises that every propositional formula can be rewritten as an equivalent formula where only the connectives **or** and **not()** appear. Consider now the operator **nand** defined by:

$$A \text{ nand } B ::= \text{not}(A \text{ and } B).$$

Prove that every propositional formula can be rewritten as an equivalent formula where only the connective **nand** appears.

Exercise 3.5 (cf. Problem 3.28)

Express each of the following statements using quantifiers, logical connectives, and/or the following predicates:

- $P(x) ::= 'x \text{ is a monkey}'$
- $Q(x) ::= 'x \text{ is a 6.042 TA}'$
- $R(x) ::= 'x \text{ comes from the 23rd century}'$
- $S(x) ::= 'x \text{ likes to eat pizza}'$

where x ranges over all living things.

- (a) No monkey likes to eat pizza.

- (b) Nobody from the 23rd century dislikes eating pizza.
- (c) All 6.042 TAs are monkeys.
- (d) No 6.042 TA comes from the 23rd century.
- (e) Does part (d) follow from parts (a), (b) and (c)? If so, give a proof. If not, give a counterexample.
- (f) Translate into English: $\forall x . (R(x) \text{ or } S(x) \text{ implies } Q(x))$
- (g) Translate into English:

$$\exists x . (R(x) \text{ and not}(Q)(x)) \text{ implies } \forall x . (P(x) \text{ implies } S(x))$$

Exercise 3.6 (from the classroom test of 3 October 2018)

Find a counter-model for the following predicate formula:

$$(\exists x . \forall y . (P(x) \text{ implies } Q(y))) \text{ implies } (\forall x . (P(x) \text{ implies } \exists y . Q(y))) .$$